# WICAT <br> Multi-user Control System (WMCS) LINK 

Programmer Reference Manual

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The Purpose of This Manual
The manual explains how to use the LINK Program for program development.

## Reader's Guide to MCS Publications

Instructions: Determine the audience to which you belong and then read only the publications at an arrowheac. Dotted arrowheads indicate optional reading.


| Volume 1 | Volume 2 | Volume 3 |
| :--- | :--- | :--- | Volume 4

An introductory user manual is a tutorial introduction to a software product. In other words, use of the product is explained in a step-by-step, user-friendly format that walks you through same fundamental aspects of the product. When you complete the introductory manual, you have the experiential basis for understanding the products's user reference manual.

User reference manuals are written for those who are new to the product (but who have read the introductory user manual) as well as for the experienced user.

System manager reference manuals contain information for those who perform administrative tasks associated with routine system operation.

Programmer reference manuals are written for programmers who at least understand programming fundamentals.

## WMCS LINK Programmer Reference Manual

## Table Of Contents

CHAPTER I OVERVIEW
1.1 ..... 1-1
THE ROLE OF THE LINKER1.1 .11.1 .2
1.2 FUNCTIONS ..... 1-4Program Modules1-1
Compilation And Assembly ..... 1-4
1.2.1 Executable-image File ..... 1-4
Cross Reference Map File 1.2 .2 ..... 1-4
SYMBOLS AND REFERENCES
2.1 DEFINITION ..... 2-1
2.2 TYPES ..... 2-2
2.2.1 Local Symbols ..... 2-2
2.2.2 Global Symbols ..... 2-3
CHAPTER 3 OBJECT MODULES
3.1 DEFINITION ..... 3-1
3.2 RECORD TYPES ..... 3-5
3.2 .1 Delimiter ..... 3-5
3.2 .2 code ..... 3-6
3.2.3 Constants/labels ..... 3-6
3.2.4 Global Symbols ..... 3-6
CHAPTER 4 IMAGE CREATION
4.1 BIT MAP ..... 4-5
4.1.1 Disk Memory Correspondence ..... 4-6
4.1 .2 Initial Memory Allocation ..... 4-6
4.1 .3 Shared Memory Page ..... 4-7
CHAPTER 5 LIBRARIES
5.1 DEFINING A LIBRARY ..... 5-1
5.2 PRELINK ..... 5-2
5.3 REFERENCING LIBRARIES ..... 5-3
5.3.1 Default Library ..... 5-3
5.3 .2 Other Libraries ..... 5-3
APPENDIX A TROUBLESHOOTING
A. 1 LINK DIAGNOSTIC MESSAGES ..... -1
A. 2 MISCELLANEOUS ERRORS ..... A-1
A. 2.1 File Empty Or Does Not Exist ..... A-2
A. 2.2 Invalid Symbol Address ..... A-2
APPENDIX B EXAMPLE OF PROGRAM DEVELOPMENT
APPENDIX C RELOCATABLE MODULE FORMAT
APPENDIX D CROSS REFERENCE MAP LISTING
APPENDIX E BIT MAPS AND RECORDS

TYPOGRAPHICAL CONVENTIONS USED IN THIS MANUAL

Uppercase letters within text indicate sample command line character strings, file designations, diagnostic messages, and reports. Such samples are in uppercase when run into the text so that you can distinguish what would be typed (or what would appear) on the screen. Examples set off from the text are in lowercase.

Bold facing indicates what you should type onto the screen, i.e., whereas uppercase characters indicate sample character strings, etc., that are part of an example, bold faced characters indicate what you must type as part of a procedure.

Square brackets, [], indicate a function key, the name of which appears in uppercase within the brackets, e.g.., [RETRN], [CTRL], etc.

Underlining is used for emphasis.

## CHAPTER 1

OVERVIEW

The WICAT Multi-user Control System (WMCS) LINK Program is a programming tool used to prepare the output of a compiler or assembler so that that output can be executed by the computer's hardware. Therefore, LINK produces a file that is an executable image.

The executable image has sharable and nonsharable segments. Sharable segments can be shared by more than one process.

LINK also creates a listing of symbols and addresses useful for cross referencing.

The description of the LINK Command, in the WICAT Multi-user Control System (WMCS) User Reference Manual, tells you how to execute LINK.

### 1.1 THE ROLE OF THE LINKER

The object modules created by language translators (including assemblers) are nonexecutable. External symbol references are unresolved. Unresolved symbol references include run-time modules required by high level languages as well as user-specified external declarations. The linker binds the symbol reference with the symbol definition and creates a memory image that can be read into memory and executed. Without a linker, modular programming is impossible, and language translators would be much more complex.

### 1.1.1 Program Modules

Modular programming, the process of combining separately compiled or assembled modules into an executable image, simplifies and enhances program development in the following ways:

1. Smaller modules are easier to write and maintain
because it is easier to find and fix errors in a smaller module.
2. Many modules are applicable to more than one task, and modules that have been used and tested in other programs can be included in new applications. This speeds program development.
3. Complicated programming requirements are simplified by breaking the requirements into smaller tasks, each of which can be allocated to a different person for implementation and debugging.
4. In same cases, the best language for one module may not be best for another part of the program. Modular programming allows modules written in different programming languages to be combined in a single application. Thus the programmer can use the language most suited to the application.

Figure 1.1 illustrates the role of the linker in modular programming wherein TIME is a computer program that interacts with the user to set the system clock on a WICAT system. The interaction is achieved through a main program written in WICAT Pascal. The clock is actually set by an assembly language program written in WICAT Assembler.


Figure 1.1. Modular Programming

```
1.1.2 Compilation And Assembly
LINK simplifies the work done by language translators because it assigns code to executable pages.
```


### 1.2 LINK FUNCTIONS

LINK produces two files:

1. An executable-image file (EXE is the file extension for this file).
2. A cross reference-map file ( $M C R$ is the file extension for this file).

These files are discussed in the following sections.

### 1.2.1 Executable-image File

An executable-image file is file type number 1. The record size of the file is 1024 bytes and the first record of the image file contains bit maps, each of which is 256 bytes long. These maps represent logical address space. Each of the bits in a map is a semaphore for a 1-Kbyte segment of address space. The first map defines the correspondence between the records on the disk and their ultimate location in memory. The second map defines the initial memory allocation. The third map indicates which pages can be shared.

The fourth map is not used on a WICAT computer.

### 1.2.2 Cross Reference Map File

This file is a listing file that specifies each module name and starting address.

## CHAPTER 2

## SYMBOLS AND REFERENCES

LINK's primary responsibility is to resolve symbolic references between modules.

### 2.1 DEFINITION

A symbol is an identifying label or name associated with one or more program statements or data area. A reference is the use of a symbol in a program statement of data definition.

Figure 2.1 illustrates a sample program called TEST. The program is written in Pascal and consists of one main routine and one subroutine.

## SYMBOLS AND REFERENCES

```
PROGRAM TEST;
VAR I : INTEGER;
PROCEDURE INCREMENT;
BEGIN
    I := I + 1;
END;
BEGIN
I := 0;
    INCREMENT;
END.
```

Fig 2.1 Sample Program
These are the symbols used in TEST:
TEST The name of the main routine
INCREMENT The name of the subroutine
I The name of a data area or variable

These are the references used in TEST:
line 7 I is referenced twice
line $11 \quad$ I is referenced once
line 12 The subroutine INCREMENT is called

### 2.2 TYPES

Symbols are one of two types: local or global. LINK treats each type differently.

### 2.2.1 Local Symbols

Local symbols can be referenced only in the routine in which the local symbols are defined.

Note that for a program such as that in figure 2.1 (where there are no local symbols) the compiler or assembler resolves all references to local symbols.

### 2.2.2 Global Symbols

Global symbols can be referenced by routines other than the routine that defines them. For example, in figure 2.1 all symbols are global because they are defined by TEST (the main routine) and are therefore global to all subroutines. LINK resolves global symbol references (see chapter 4).

Global symbols are of two types:

1. Intemally defined symbols, defined in the main program.
2. Externally defined symbols, defined in a routine that is external to the main program, and independently compiled (e.g., runtime routines used by Pascal).

## CHAPTER 3

OBJECT MODULES

### 3.1 DEFINITION

The output (or object) file for the compiler or assembler is the input file for LINK. An object file consists of modules. In the case of the Pascal compiler, a discrete module is created for each procedure in the program along with one for the main program itself. An additional module is created for global variables. In chapter 2, figure 2.1, we examined a simple Pascal program. Figure 3.1 illustrates the modules created by the compiler.


Fig. 3.1 Modules Created By PASCAL Compiler For TEST Program

A module can contain several records. An object record has the following hexadecimal format displayed in bytes:


Fig. 3.2 Object Record Format
This is an explanation of the terms appearing in figure 3.2:
Type Type of object record described in 3.2, a label.

Byte Count Number of bytes in the record following the record type (not including the byte
count itself).
Relative Address Relative address of the code bytes in the module.

Code Byte In the case of a code type record, the code bytes contain the actual hexadecimal code for the module; in the case of a symbol, the code bytes contain a name.

The checksum is the one's complement of the sum of the bytes.

Diagrams showing the format of each type of record are found in appendix D.

Figure 3.3 shows the hexadecimal doject file created by the Pascal compiler for program TEST (shown in figure 2.1). Note, the three modules relating to figure 3.1.


Fig. 3.3 Object File For TEST Program
Each record begins with a record label, 'R', followed by an integer between 0 and 9 or one of the alphabet characters 'I' or ' $P$ '. References to global symbols (subroutines, variables, etc.) must be made using the M68000 absolute long addressing mode. In the object file, using figure 3.3 as an illustration, all references to internally and externally defined global symbols (refer to chapter 2) are enclosed in a box. Note, all references to internally defined global symbols have an offset into the global module while all references to externally defined global symbols have an address value of 0 . These global references can be better understood by studying a listing of the assembly code generated by the compiler for TEST. This listing is illustrated in figure 3.4 .

## OBJECT MODULES

| WICAT Pascal Version 1.3 |  |  |  |
| :---: | :---: | :---: | :---: |
| ;procedure: increment( |  |  | 1) |
| 0000 | 4E56 | link | a6,\#-4 |
|  | FFFC |  |  |
| 0004 | 200D | move. 1 | a5,-(a6) |
| 0006 | 5279 | addq.w | \#1,74 |
|  | 0000 |  | global |
|  | 004A |  |  |
| 000C | 2A5E | move. 1 | $(\mathrm{a} 6)+, \mathrm{a} 5$ |
| O00E | 4E5E | unlk | a6 |
| 0010 | 4E75 | rts |  |
| ;procedure: test ( 0) |  |  |  |
| 0000 | 2F3C | move. 1 | \#6,-(sp) |
|  | 0000 |  | global |
|  | 0006 |  |  |
| 0006 | 2F3C | move. 1 | \#40,-(sp)global |
|  | 0000 |  |  |
|  | 0028 |  |  |
| 000C | 4EB9 | jsr | $\begin{aligned} & \operatorname{xxxxxx} \\ & \text { RRRO00 } \end{aligned}$ |
|  | 0000 |  |  |
|  | 0000 |  |  |
| 0012 | 4279 | clr.w | 74 global |
|  | 0000 |  |  |
|  | 004A |  |  |
| 0018 | 2A4E | move. 1 | a6, a5 |
| 001A | 4EB9 | jsr | xxxxxx INCREMENTI |
|  | 0000 |  |  |
|  | 0000 |  |  |
| 0020 | 4 EB 9 | jsr | xxxxxx RRROO1 |
|  | 0000 |  |  |
|  | 0000 |  |  |

Fig. 3.4 Assembly Code Generated By Compiler For TEST Program Note the following:

1. Global reference for I following address 0006 in module: INCREMENT
2. Global references for $I$ in module: TEST, following address 0012.
3. External reference for runtime routine RRR000 following address $C$.
4. External reference for INCREMENT module following address 1A.
5. External reference of runtime routine RRROO1 following address 20.

### 3.2 RECORD TYPES

There are four broad categories into which all object module records can be classified:

1. Delimiter.
2. Code.
3. Constant/label definitions.
4. References to global symbols.

### 3.2.1 Delimiter

There are six record types in this category, as defined by the second character in the type field:

RP Delimits a sharable segment; pure storage.
RI Delimits a nonsharable segment; impure storage.
RO Identifies the main program. The program name is coded in the code bytes. This record delimits the sequence of records pertaining to the main program.

Rl Identifies a module or subroutine. The module name is coded in the code bytes. This record delimits the sequence of records pertaining to a module.

R8 This record marks the end of a module. The address of the first free location past the end of the module. LINK uses this offset from the start of the module to determine where it can start loading the next module. This record contains the amount of space used by the module.

R9 This record marks the end of the object file. An object file can contain several modules.

### 3.2.2 Code

An R2 label indicates that the record contains program code.

### 3.2.3 Constants/labels

There are four types of object records found in this category:

R3 Identifies a local constant. The name is coded in the code bytes. The relative address field contains the value of the constant.

R4 Identifies a local label. The name is coded in the code bytes. The relative address field contains an offset pointing to the labeled location.

R5 Identifies a global constant. The name is coded in the code bytes.

R6 Identifies a global label. The name is coded in the code bytes.

### 3.2.4 Global Symbols

Record type $R 7$ identifies a reference to any global symbol. The relative address field is the location (within the module) of the reference to the global symbol.

## CHAPTER 4

## IMAGE CREATION

LINK reads relocatable object modules, resolves refererices between modules, and creates an image file that the operating system can load and execute.

As relocatable modules are read in, they are placed in a virtual address space that resembles the address space the program uses during execution. By convention, this address space is divided in half with the lower half reserved for program code and constants. The upper half is reserved for the stack and other variables. The purpose of this convention is to allow the non-changing or pure portions of programs to be shared among several users. See figure 4.1 .

IMAGE CREATION


Fig. 4.1 Memory Map Of The 2-Mbyte Virtual Address Space
Space is allocated consecutively in these two segments. Allocation in each area is controlled by a base pointer that points to the lowest unallocated memory location. The pointer into the lower half of the address space is called Pure-Base and has an initial value of 1000 Hex. The pointer into the upper half of the address space is called Impure-Base and has an initial value of 100000 Hex .

As the relocatable modules are read in, they are placed into the virtual address space relative to one of these two bases. The base pointer chosen depends on the 'RP' or 'RI' type record that must be the first record in a module. If the record is type 'RP' the module goes into the lower half of the address space. If the record is type 'RI' the module goes into the upper half of the address space.

The end of a module is delimited by an 'R8' type of a record. The address field of an 'R8' type record contains a count of the number of bytes of memory used by the module. When a module has been completely read in, the base pointer is incremented by the value of the address
field in the R 8 record. Thus the pointer once again points to the battam of free memory.

Between the start and end of a module, three classes of records may be encountered. The first class causes data to be entered into the virtual memory address space. The second class causes symbols to become defined. The third class marks a reference to a symbol.

There is only one instance of the first record class. This is the 'R2' type of record. This record causes bytes of data to be stored in the virtual memory space.

The second and third record classes affect the linker symbol table that is defined by the series of Pascal statements in figure 4.2.


Fig. 4.2 PASCAL Definition Of Linker Symbol Table
Pictorially, the symbol table corresponds to the diagram in figure 4.3.


Fig. 4.3 Pictorial Representation of LINK Symbol

The second record class consists of types 'RQ', 'Rl', 'R5', and 'R6'. These result in a symbol table entry being built if one does not yet exist and cause the value field of the symbol table entry to become defined.

The third record class consists of only record type 'R7'. An occurence of a type 'R7' record causes a symbol table entry to be built if one does not yet exist. An entry is then made into a list of locations of undefined symbols that is associated with the symbol table entry.

After all of the relocatable modules are read in, LINK procedes to add the value field of each symbol table entry to all of the locations where it is referenced as specified in the associated list of references to undefined symbols.

After the undefined references have been resolved, the executable image file is created.

### 4.1 BIT MAP

The executable image file has a record length of 1024 bytes. The first record of this file contains bit maps each of which is 256 bytes long. These bit maps are:
l. Disk Memory Correspondence.
2. Initial Memory Allocation.
3. Shared Memory Page.
4. NOT USED.

These maps are arranged as shown in figure 4.4.


Fig. 4.4 Image File Format

### 4.1.1 Disk Memory Correspondence

The first bit map defines the relationship between the 2-Mbyte logical address space and the records of the executable image file. Each bit of this map corresponds to 1024 bytes of memory. If the bit is set, the next record from the executable image file is read into memory at the location corresponding to the bit position.

### 4.1.2 Initial Memory Allocation

The second bit map indicates which pages of memory should be allocated to the new process. Again, each bit corresponds to 1024 bytes in the 2-Mbyte address space. The operating system allocates memory to 4096-byte pages. Therefore, the setting of any of the bits on the page causes the full 4096 bytes to be allocated.

### 4.1.3 Shared Memory Page

The third bit map is the memory-protection bit map. The setting of any of the bits within the range corresponding to a 4096 -byte page causes the full 4096 bytes to be write protected by the operating system. This process makes the page sharable because the contents cannot be changed.

There is space for a fourth bit map that is unused at this time and is assumed to be zero.

## CHAPTER 5

## LIBRARIES

LINK uses the content of a library definition file to resolve undefined references after all the user-specified files have been input.

### 5.1 DEFINING A LIBRARY

To define a library, you create a file containing a list of equivalences. An entry in this file is a symbolic name that can be referenced followed by the filename of an object module that defines the symbolic name.

You can create this file as a normal text file by using the VEW Program (read the Virtual Editing Window (VEW) User Reference Manual). The standard extension for a library definition file is .DEF. When referencing a file in an entry, you should enter the filename in uppercase and include the file extension.

A line that begins with an ampersand, @, indicates indirection through another library definition file. In this case, LINK responds as though the contents of the other file had been inserted at this point.

Figure 5.1 is an example of a library definition file.

| $\underset{12345678901234567890}{2}$ | $\begin{array}{ccr} 3 & 4 & 5 \\ 123456789012345678901234567890 \end{array}$ |
| :---: | :---: |
| RRROOO | _DSO/PRIT IB/INIT.MRL |
| RRROO1 | _DS0/PRILIB/EXIT.MRL |
| _PUTSTRING | _DSO/PRILIB.SVC/PUTSTR.MRL |
| @_DSO/RLIB/ANOTHER.DE |  |

Fig. 5.1 Library Definition File

### 5.2 PRELINK

When LINK is ready to use a file, it must translate the name of the file to get the File Control Block (FCB) number assigned to the file. You can often save time during the linking process by performing this translation as a separate step. The PRELINK Program does this. Execute PRELINK by typing either of the following command-line character strings and then striking [REIRN] :
> prelink linklib.def
or
> prelink _ds0/syslib/linklib.def
PRELINK assumes the file has a .DEF extension if no file extension is specified in the PRELINK command-line character string.

PRELINK adds another column to the list of equivalences in the .DEF file. This column consists of the FCB numbers of the specified files. See figure 5.2 for an example of a library definition file after processing by PRELINK.

| $\begin{array}{cr} 1 & 2 \\ 12345678901234567890 \end{array}$ | $\begin{array}{ccr} 3 & 4 & 5 \\ 123456789012345678901234567890 \end{array}$ | $12345678901234567890^{6}$ |
| :---: | :---: | :---: |
| RRROOO | DSO/PRILIB/INIT.MRL | _DS0//\#96.1 |
| RRROO1 | DSO/PRILIB/EXIT.MRL | _DSO//\#97.3 |
| _PUTSTRING | DSS/PRILIB.SVC/PUTSTR.MRL | _DS0//\#106.1 |

@_DMO/RLIB/ANOTHER.DEF
Fig. 5.2 Library Definition File After Processing By Prelink

Note: PRELINK creates a new version of the library definition file in the same directory as the original file. To
conserve disk space, you may want to purge old library definition files after running PRELINK.

### 5.3 REFERENCING LIBRARIES

### 5.3.1 Default Library

LINK always references the file LINKLIB.DEF, found in directory /SYSLIB/ on the system disk. During initialization, LINK reads the contents of this file to make the names of library routines available to programs.

### 5.3.2 Other Libraries

To reference another library definition file, use the :LIBRARY switch on the command line. For example:
> link one,two,three :library=test
The foregoing command tells LINK to link files ONE, TWO, and THREE and reference TEST.DEF as a library file.

## APPENDIX A

TROBLLESHOOTING

## A. 1 LINK DIAGNOSTIC MESSAGES

These are the diagnostic message associated with LINK: Display Undefined Symbols

Checksum Error In Reading Relocatable Module
Unexpected End of File While Reading Relocatable Module
Address Of Undefined Reference Accesses Unallocated Disk Page

Open Error (WMCS diagnostic message number)
Read Error (WMCS diagnostic message number)
Write Error (WMCS diagnostic message number)
Close Error (WMCS diagnostic message number)
Delete Error (WMCS diagnostic message number)

## A. 2 MISCELLANEOUS ERRORS

Certain conditions lead to errors that do not give the preceding error messages. Some of these conditions are discussed in this section.

## A.2.1 File Empty Or Does Not Exist

If a specified file is empty or does not exist, the message 'Unexpected end of file encountered while reading' is displayed followed by the name of the file. The solution is to use the TYPE Command to display the contents of the file. Be certain to specify the filename given in the diagnostic message.

## A.2.2 Invalid Symbol Address

LINK does not check the validity of symbol addresses. If for some reason, a symbol has an offset beyond the end of its module, this error can happen to be caught during the process of adding the actual address to the location of an undefined reference. The solution is to examine the source modules for unreasonable address offsets.

## APPENDIX B <br> EXAMPLE OF PROGRAM DEVELOPMENT

The diagram on the next page shows how files produced by the assembler and higher-level language compilers may be linked to form an executable image. A number of relocatable code (.MRL) files can be linked in one operation.

EXAMPLE OF PROGRAM DEVELOPMENT


## APPENDIX C

RELOCATABLE MODULE FORMAT

The relocatable source file contains a series of modules terminated by an R9 record. A module is a series of record types R2-R7 preceded either by an RO or an Rl type record and terminated by an R8 type record.


TYPE COUNT ADDRESS


The byte count is the number of bytes not including the byte count itself.


TYPE COUNT ADDRESS SUM


The checksum byte is the ones complement of the sum of the bytes. This form of checksum was chosen to be compatible with Motorola.

RELOCATABLE MODULE FORMAT

| REC TYPE | BYTE COUNT ADDRESS | PROGRAM NAME | $\begin{gathered} \text { CHECK } \\ \text { SUM } \end{gathered}$ |
| :---: | :---: | :---: | :---: |
| \| RO | BC | Al | A2 | A3 | A4 | \#\# | \#\# | \#\# | \#\# | ....| \#\# | \#\# | CS | |  |  |  |
| Delimits the sequence of records pertaining to the main program. Otherwise, this type is treated the same as type 6. |  |  |  |
| REC | BYTE |  | CHECK |
| TYPE | COUNT ADDRESS | MODULE NAME | SUM |
| \| R1 | BC | Al | A2 | A3 | A4 | \#\# | \#\# | \#\# | \#\# | .... | \#\# | \#\# | CS | |  |  |  |

Delimits the sequence of records pertaining to a subroutine. Otherwise, this type is treated the same as type 6.


Local constant:
The label is defined to have the value of the address field.


Local label:
The label is defined to have the value of the address field plus the relocation constant.


Global constant:
The label is defined to have the value of the address field.


## Global label:

The label is defined to have the value of the address field plus the relocation constant.

| REC | BYTE |  | CHECK |
| :---: | :---: | :---: | :---: |
| TYPE | COUNT ADDRESS | LABEL | SUM |
| R7 | BC \| Al | A2 | \#\# \| \#\# | CS |

The address field of this record is the location of the reference to the undefined label.

| REC <br> TYPE COUNE <br> COUNT ADDRESS |
| :--- |
| R8 \| BC | AI | A2 | A3 | A4 | CS |CHECK <br> SUM |

The foregoing kind of record marks the end of a module. This -record contains the amount of space used by the module. The address field from the record is added to the current location pointer to form the address of the first free location past the end of the module. The address of the first free location past the end of the module then becomes the new location pointer.

RELOCATABLE MODULE FORMAT

| REC |
| :--- |
| TYPE CYTE |
| COUNI ADDRESS | | CHECK |
| :---: |
| SUM |

The foregoing kind of record marks the end of file. The address field is undefined.

## APPENDIX D

CROSS REFERENCE MAP LISTING

The following pages display a cross reference map listing from the sample program TEST.

WICAT Link Editor - Version 2.2
Linker map :

File : TEST.MRL .START. INCREMENTI TESTO Module \$PASCAL-GLOBAL\$-2 Module

File : _DIO/PRTLIB/EXIT.MRL RRROO1

Label
File : _DIO/PRILIB/INIT.MRL RRROOO

Label
File : _DIO/PRILIB/RIGIOBAL.MRL RIGTOBAL HEAP Module Module

File : _DIO/PRIIIB/RENRITE.MRL RRROIl Label

File : _DIO/PRILIB/RESET.MRL RRROIO Label

File : _DIO/PRILIB/FBINIT.MRL RRROO9

Label
File : _DIO/PRILIB/SHOWIO.MRL SHOWIO

Label
File : _DIO/PRILIB/ERROR.MRL ERROR

Label
File : _DIO/PRILIB/HANG.MRL
File : _DIO/PRILIB/WHEX.MRL RRR039

File : DIO/PRILIB/PUT.MRL RRRO22

File : _DIO/PRIIIB/GET.MRL RRRO21

Label
File : _DIO/PRITIB/MBVALID.MRL RRR020 Label

|  | Base - Pure= | 1000 | Impure $=100000$ |
| :---: | :---: | :---: | :---: |
| 1000 |  | 6 | 0 |
| 1006 |  | 12 | 0 |
| 1018 |  | 26 | 0 |
| 100000 |  | 0 | 4 C |
|  | Used - Pure= | 3E | Impure $=14 \mathrm{C}$ |
|  | Base - Pure= | 103E | Impure $=10004 \mathrm{C}$ |
| 103E |  | 40 | 0 |
|  | Used - Pure= | 40 | Impure= 0 |
|  | Base - Pure= | 107E | Impure $=10004 \mathrm{C}$ |
| 107E |  | 126 | 0 |
|  | Used - Pure= | 126 | Impure $=0$ |
|  | Base - Pure= | 11A4 | Impure $=10004 \mathrm{C}$ |
| 10004C |  |  |  |
| 10005C |  | 0 | 20 |
|  | Used - Pure= | 0 | Impure $=20$ |
|  | Base - Pure= | 11A4 | Impure $=10006 \mathrm{C}$ |
| 11A4 |  | C4 | 0 |
|  | Used - Pure= | C4 | Impure $=0$ |
|  | Base - Pure= | 1268 | Impure $=10006 \mathrm{C}$ |
| 1268 |  | EA | C |
|  | Used - Pure= | EA | Impure= 0 |
|  | Base - Pure= | 1352 | Impure=10006C |
| 1352 |  | 4 C | 0 |
|  | Used - Pure= | 4 C | Impure= 0 |
|  | Base - Pure= | 139E | Impure $=10006 \mathrm{C}$ |
| 139E |  | 60 | 0 |
|  | Used - Pure= | 60 | Impure= 0 |
|  | Base - Pure= | 13FE | Impure=10006C |
| 13FE |  | 52 | 0 |
|  | Used - Pure= | 52 | Impure= 0 |
|  | Base - Pure= | 1450 | Impure $=10006 \mathrm{C}$ |
| 1450 |  | 48 | 0 |
|  | Used - Pure= | 48 | Impure= 0 |
|  | Base - Pure= | 1498 | Impure $=10006 \mathrm{C}$ |
| 1498 |  | 98 | 0 |
|  | Used - Pure= | 98 | Impure= 0 |
|  | Base - Pure= | 1530 | Impure $=10006 \mathrm{C}$ |
| 1530 |  | 17 C | 0 |
|  | Used - Pure= | 17C | Impure= 0 |
|  | Base - Pure= | 16AC | Impure $=10006 \mathrm{C}$ |
| 16AC |  | 7 C | 0 |
|  | Used - Pure= | 7 C | Impure= 0 |
|  | Base - Pure= | 1728 | Impure=10006C |
| 1728 |  | 1 CO | 0 |
|  | Used - Pure= | 1 CO | Impure= 0 |


| File : _DIO/PRILIB/CTOSE.MRL <br> ** Unable to open by FCB! |  |  | Base - Pure= | 18E8 Impure=10006C |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| RRRO19 | Label | 18E8 |  |  |  |  |
| RRR019A | Label | 18 E 8 |  | F4 |  | 0 |
|  |  |  | Used - Pure= | F4 | Impure= | 0 |
| _HEAP | Symbol | 10006C |  | 0 |  | 0 |
| Image filename = TEST.EXE |  |  |  |  |  |  |
| Total image size (all sizes g | given in | xadecima | 1) $=$ |  |  |  |
| 9DC bytes pure code. |  |  |  |  |  |  |
| 6C bytes impure code. |  |  |  |  |  |  |
| 0 bytes reserved for | dynami | pace. |  |  |  |  |
| 1000 bytes reserved for | stack. |  |  |  |  |  |

## CROSS REFERENCE MAP LISTING

WICAT Link Editor - Version 2.2
Symbol Cross Reference

| \$PASCAL-GLOBAL\$-2 | 100000 | 100E | 101A | 1020 | 102C |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| .START. | 1000 |  |  |  |  |  |  |  |
| ERROR | 13FE | 1230 | 131E | 13B0 | 13D2 | 145E | 156A | 16E6 |
|  |  | 18B0 | 19A4 |  |  |  |  |  |
| HANG | 1450 | 123 E | 132C | 1578 | 16 F 4 | 18BE | 19B2 |  |
| HEAP | 10005 C | 1150 |  |  |  |  |  |  |
| INCREMENII | 1006 | 1034 |  |  |  |  |  |  |
| RRR000 | 107E | 1026 |  |  |  |  |  |  |
| RRRO01 | 103E | 103A | 1454 |  |  |  |  |  |
| RRR009 | 1352 | 1102 | 1130 | 198A |  |  |  |  |
| RRRO10 | 1268 | 1140 |  |  |  |  |  |  |
| RRRO11 | 11A4 | 1112 |  |  |  |  |  |  |
| RRRO19 | 18E8 | 11BA | 127E |  |  |  |  |  |
| RRRO19A | 18 E 8 |  |  |  |  |  |  |  |
| RRRO20 | 1728 | 16C6 |  |  |  |  |  |  |
| RRRO21 | 16AC | 12E4 |  |  |  |  |  |  |
| RRRO22 | 1530 | 1514 |  |  |  |  |  |  |
| RRR039 | 1498 | 13C8 |  |  |  |  |  |  |
| RTGLOBAL | 10004 C | 1044 | 105A | 1118 | 1146 | 13B6 | 1408 |  |
| SHOWIO | 139E | 1238 | 1326 | 1572 | 16EE | 18B8 | 19AC |  |
| TESTO | 1018 | 1002 |  |  |  |  |  |  |
| _HEAP | 10006C | 115A |  |  |  |  |  |  |

## APPENDIX E

## BIT MAPS AND RECORDS

This appendix contains bit maps and records for the sample program TEST. These bit maps and records are a hexadecimal dump of the image file for the TEST program.

## FILE: _DIO/TOM.CDS/TEST.EXE.I



| 000005C0 | 000005D0 000005E0 000005F0

000800 OE $422800014 \mathrm{~A} 68001856 \mathrm{E} 80016 \mathrm{~B}(\mathrm{JhVh}$ 22 6E 00 OA 2009 4A 1966 FC 908944805380 nn Jf|DS 2F 2E 00 OA 2F $002 \mathrm{~F} 3 \mathrm{C} 000000023 \mathrm{~F} 280014 / . / /<$ ? ( 426742 A 72 F 3C FF FF FF FF 2F 3C FF FF FF FF BgB' $/</<$ Record

00000600 00000610 00000620 00000630 00000640 0000065 0000066 00000670
0000068 00000690 000006AO 000006B0 000006C0 000006D 000006E 000006F0 48680008486800044 E 4000004 A A8 0004 HhHhNe 57 E8 00 OD 66142 D 6E 000400 OE 4C DF 0301 Whf-nil
 000013 FE 2F 084 E b9 000013 9E 4E F9 $0000 \sim / \mathrm{N} 9 \mathrm{Ny}$ 14505252523031312020202052455752 PRRROII - REWR 495445202 D 2 D 20204352454154204641 ITE - CREAT FA 494 C 45440 D 0 A 00004 E 56000048 E 80 CO ILEDNHIg 206 E 000 E 4 A 28000 D 67082 F 084 E B9 $0000 \mathrm{~nJ}(\mathrm{~g} / \mathrm{N}$ 18 E8 $316 \mathrm{E} 0008000 \mathrm{E} 42280016226 \mathrm{E} 00 \mathrm{OA} \mathrm{hlnB}(\mathrm{n}$ 20094 A 1966 FC 9089448053802 F 2 E 000 A JfldS/. 2F 002 F 3C $000000013 \mathrm{~F} 28001442674868 / /<$ ? (BgHh 0008486800044 E 4000024 A A8 $0004663 \mathrm{CHhN@J}$ (f< 50 E8 00 OD 422800024228000350 E 80000 PhB (B(Ph 217 C FF FF FF FF 0010217 C 0000000100 lC ! ! ! 2F 084 E B9 $000016 \mathrm{AC} 2 \mathrm{D} 6 \mathrm{E} 0004000 \mathrm{E} 4 \mathrm{C} D \mathrm{DF} / \mathrm{Na},-\mathrm{nL}$
|
| 0000070
| 0000071
| 0000072
| 0000073
| 0000074
| 0000075
| 0000076
| 00000770
| 00000780
| 0000079
| 000007A0
| 000007B0
| 000007C
|000007D
| 000007E
| 000007F

| $\mid$ | 00000800 |
| :--- | :--- |
| $\mid$ | 00000810 |
| $\mid$ | 00000820 |
| 1 | 00000830 |
| $\mid$ | 00000840 |
| $\mid$ | 00000850 |
| $\mid$ | 00000870 |
| $\mid$ | 00000880 |
| $\mid$ | 00000890 |

FF F8 48 E7 FF FC 20 7C $0010004 \mathrm{C} 20504286 \mathrm{xHg} \mid \mathrm{LL}$ PB 226 E 0008 4A 196704528660 F8 2F 280008 " $\mathrm{nJgR} \times \mathrm{x} /($
 $00082 \mathrm{~F} 06486 \mathrm{EFFF8} 486 \mathrm{EFF} \mathrm{FC} 4 \mathrm{E} 400005 / \mathrm{HnxHn} \mid \mathrm{N@}$
 2000 4E F9 000010 3E 48 7A 00 OA 4E B9 $0000 \quad \mathrm{Ny}>\mathrm{HzN} 9$ 13 FE 60 FE 45584543555449 4E 4720224 A ~~ ${ }^{\text {~ EXECUTING "J }}$ $554 D 5020544 F 204845524522204 C 4 F 4 F$ UMP TO HERE" LOO $5020494 E 2050415343$ 41 4C 2052554 E 54 P IN PASCAL $\mathrm{P}^{\mathrm{r}-\pi}$ 494 D 45 OD OA 000000 4E 56 FF EC 48 E7 FO CO IMENVIHgp@

000008AO 000008B0 000008C0 000008D0 000008EO 000008F0

00000900 00000910 00000920 00000930 00000940 00000950 00000960 00000970 00000980
| 00000990
| 000009AO
| 000009BO
| 000009C0
| 000009D0
| 000009E0
| 000009F0

206 E 00 OE 7014224 E 13 3C $0020538066 \mathrm{~F} 8 \mathrm{np} \mathrm{n}^{\mathrm{N}}<\mathrm{N}<\mathrm{Sfx}$ 428324 2E 00 OA 22 4E 2202 E8 8A $02810000 \mathrm{~B} \$ . \mathrm{N}^{\mathrm{N} \mathrm{N}^{\prime} \mathrm{h}}$ 00 OF OC $0100096 F 025 E 010601003013010^{\wedge} 0$ 5283 4A 8266 E2 32 2E 000848 Cl 6A 022203 RJfb2. $\mathrm{HAj} \mathrm{J}^{n}$ B2 83 6C OC 22 4E 200113 3C 00 2A 53806 E F8 $21^{\prime \prime} \mathrm{N}$ <*Snx 4A 816728 B6 81 6A $08117 C 002000206004 \mathrm{Jg}(6 j)$

11590020117 C 00010000117 C 00010001 Y || 2F 084 E B9 00001530538160 D 42 D 6 E 0004 /N90S`T-n 00 OE 4C DF 03 OF 4E 5E DF FC 000000 OA 4E 75 L_NN_|Nu 4E 56 FF EE 2F 08206 E 0008 4A \(28000067 \mathrm{lE} \mathrm{NVn/} \mathrm{~nJ}(\mathrm{~g}\) 6100009252 A8 \(00104228000042280001 \mathrm{aR}(\mathrm{B}(\mathrm{B})\) 2D 6E 00040008205 F 4 E 5 E 58 8F \(4 \mathrm{E} 75487 \mathrm{~A}-\mathrm{n}\) _ \(\mathrm{N}^{\wedge} \mathrm{XNuHz}\) 00 1C 600448 7A 00 4C 4E B9 000013 FE 2 F 08 `HzIN9~/ 4E B9 000013 9E 4E F9 0000145052525230 N9NyPRRRO

 41204953204 E 4 F 542056414 C 4944 OD OA A IS NOT VALID 00005252523032322020 2D 2D 20505554 RRRO22 - PUT 20 2D 2D 2020575249544520464149 4C 45 —— WRITE FAILE 44 OD OA 0048 E7 806042 AE FF FC 4A 2800 OC DHg`B.|J( 6708 2D 7C 00000003 FF FC 4A 6800186700 g-l|Jhg 00 8A 42 6E FF EE 3D 680018 FF F0 42 6E FF F2 Bnn=hpBnr
| 00000A00
00000A10
00000A20
00000A30
00000A40
00000A50 00000A60 00000A70 00000A80 00000A90 00000AAO 00000ABO 00000ACD 00000ADO 00000AEO 00000AFO

3D 680014 FF F 42028001090 A8 00 lC 5280 =ht ( R BO AE FF EE 5C EE FF FA CO E8 001443 E8 $00200 . \mathrm{n} \backslash \mathrm{nz} @ h C h$ 2449 D5 CO 20 2E FF F2 600214 D9 51 C8 FF FC $\$ 1 \mathrm{C} @$.r`YQH| 4A 2E FF FA 670000 3E 2F 2800082 F 2800 1C J. $2 \mathrm{~g} / \mathrm{/} /(/$ ( $2 F 2 E F F F C 2 F 3 C$ FF FF FF FF $48680020202 \mathrm{E} / .1 /<\mathrm{Hh}$. FF F2 Dl 97 2F 2E FF EE 48 6E FF F6 $48680004 \mathrm{rQ} /$. $\mathrm{nH} \mathrm{H} v \mathrm{VHh}$ 4E 400005 4A A8 00046600 FE FA 20 2E FF EE N@J ( $£ \sim 2$. DI A8 00 1C 4C DF 06014 E 75 2F 2800082 F 28 Q(L_Nu/(/) 00102 F 2 EFF FC 2 F 3 C FF FF FF FF $48680020 / .1 /<\mathrm{Hh}$ 2F 3C 0000000148 6E FF F6 $486800044 \mathrm{E} 40 /<\mathrm{HnvHhN}$ @ 0005 4A A8 000467 CC 6000 FE BA 4E $560000 \mathrm{~J}\left(\mathrm{gL}{ }^{\wedge} \sim: N V\right.$ 2F 0820 6E 0008 4A 2800026624 4A $280000 / \mathrm{nJ}(\mathrm{f} \$ \mathrm{~J}($ 66082 F 084 E B9 $000017284228000052 \mathrm{~A} 8 \mathrm{f} / \mathrm{N} 9(\mathrm{~B}(\mathrm{R})$ 00102 D 6 E 00040008205 F 4 E 5 E 58 8F 4E $75-\mathrm{n}$ _ N ^ XNu 48 7A 0016 4E B9 000013 FE 2 F 08 4E B9 $0000 \mathrm{HzN} 9 \sim / \mathrm{N} 9$ 13 9E 4E F9 $0000145052525230323120 \quad 20$ NyPRRRO21

00000в00 00000B10 00000B20 00000B30 00000B40 00000B50 00000B60 00000B70

2D 2D 2047455420 2D 2D $2020454 F 462057$ — GET — EOF W 4153205452554520505249 4F 522054 4F AS TRUE PRIOR TO $204745540 D O A 00004 \mathrm{E} 56 \mathrm{FF}$ EE 2F 0820 6E GEINVn/ n 0008 4A 28000066000048 4A $2800026600 \mathrm{~J}(\mathrm{fHJ}(\mathrm{f}$ $00406100004 A 4 A 28000 C 673442280003$ @aJJ (g4B ( OC 2800 OD 002067240 C 2800 OA $002067 \mathrm{lC}(\mathrm{g} \$ \mathrm{c} \mathrm{g}$ OC 2800 OB 00206714 OC 2800 OC 002067 OC (g (g 4A $280020660 \mathrm{~A} 117 \mathrm{C} 0020002050 \mathrm{E} 80003 \mathrm{~J}(\mathrm{fl} \mathrm{Ph}$



Record| !

[^0]
| 00001700
00000000000000000000000000000000 0000171000000000000000000000000000000000 0000172000000000000000000000000000000000 0000173000000000000000000000000000000000 0000174000000000000000000000000000000000 0000175000000000000000000000000000000000 0000176000000000000000000000000000000000 0000177000000000000000000000000000000000 0000178000000000000000000000000000000000 0000179000000000000000000000000000000000 000017AO 00000000000000000000000000000000 $000017 B 000000000000000000000000000000000$ 000017CO 00000000000000000000000000000000 000017D0 00000000000000000000000000000000 OOOO17EO 00000000000000000000000000000000 000017FO 00000000000000000000000000000000

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Briefly descrice examples, illustrations, or information that you think should be acced to this manual.
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What would you delete from the manual and why?
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$\qquad$
$\qquad$
What areas need greater emphasis?
$\qquad$
$\qquad$
$\qquad$
List any terms or symbols used incorrectly.
$\qquad$
$\qquad$
$\qquad$

First Fold



[^0]:    | 0000140000001002000010 OE 000010 1A 00001020 | $00001410000010260000102 \mathrm{C} 000010340000103 A \&, 4$ :

